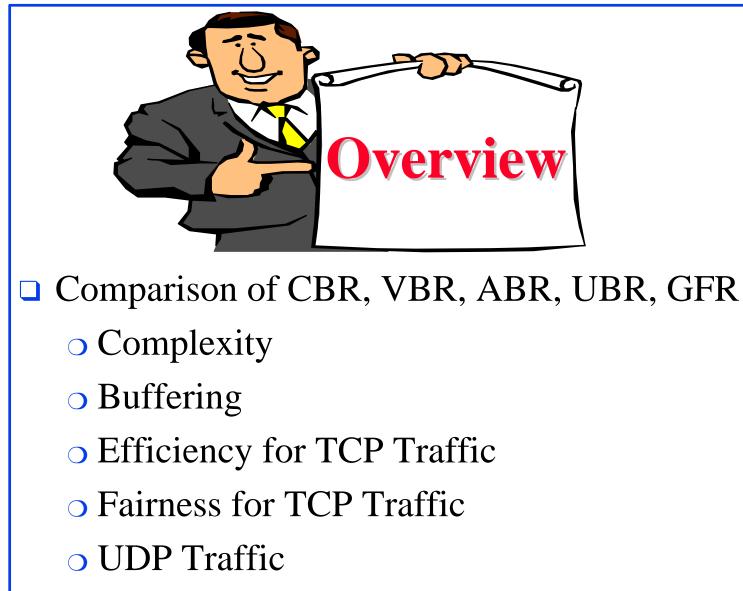
<b>A Comparison of</b>
<b>ATM Service</b>
Categories
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• Differentiated Services

#### Issues

- Services: CBR, VBR, ABR (with MCR), UBR (no MCR), GFR (with MCR)
- UBR with MCR has characteristics in-between UBR and GFR
- □ VBR  $\Rightarrow$  nrt-VBR (except in voice discussion)
- Metrics: Cost/Complexity, Performance (throughput, buffering, fairness)
- Applications: Data (TCP or UDP), Voice, Differentiated Services
- Configurations: Backbone ATM, end-to-end ATM
- Note: No absolute answers. Only points for a debate. The Ohio State University
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## Complexity

- □ Note: Service categories are listed best first.
- □ CAC (Provisioning): UBR, CBR, ABR, GFR, VBR
- Delicing: UBR, CBR, VBR, GFR, ABR
- Meeting Service Guarantees in Switches (Resource Allocation algorithm): CBR, nrt-VBR, rt-VBR, UBR (need frame boundaries), GFR, ABR
- VC Aggregation: CBR, UBR, ABR, GFR (different frame sizes), VBR
- Queueing (# of queues for n VCs): UBR (1), CBR/VBR/ABR/GFR (n) The Ohio State University

## **Complexity (Cont)**

#### Complexity of Implementation (Switch cost, NIC cost): CBR, UBR, VBR, ABR, GFR

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#### **Switch Buffering**

- **CBR:** Almost no buffering
- □ ABR: Low buffering
- □ VBR/GFR/UBR: High buffering

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#### **Router or End-system Buffering**

- Depends on the type of traffic
- □ UBR, GFR, VBR: Traffic immediately enters the ATM network ⇒ Low buffering
- CBR: Queues depend upon peak traffic rate and PCR
   ABR:
  - Queues in the end systems or routers
  - Ack regulation schemes can control required buffering for TCP

#### **Use of Extra Router Buffering**

- ABR/CBR: Routers can buffer when the backbone network is congested.
   Waiting is generally better than loss.
- GFR/VBR/UBR: Router does not know about network congestion. Extra memory does not help.

## **Bursty TCP Traffic: Bandwidth Utilization**

- $\Box$  High Utilization  $\Rightarrow$  Less idle time
- ABR: Any available bandwidth is immediately allocated
- □ GFR/UBR/VBR: Higher burstiness
   ⇒ More queues/loss and More idle times
- □ CBR: Not suited for bursty traffic

## **Bursty TCP Traffic - Fairness**

Configuration I:

ATM backbone  $\Rightarrow$  VCs between Routers

 $\Rightarrow$  Each VC carries multiple TCP flows

- ABR: Most losses in the router not in switches
   ⇒ Key factor is the fairness in the router
   ⇒ Proper RED can make it fair
- **CBR:** Queues in routers (as in ABR)
- □ VBR/GFR/UBR:

Not fair since most losses in ATM switches. Fair buffer allocation (FBA) can ensure fairness among VCs but not among flows in the same VC.

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## **Bursty TCP Traffic - Fairness**

Configuration II:

- ATM end-to-end  $\Rightarrow$  1 VC per TCP flow
- □ ABR: No losses
- **CBR:** No losses
- GFR: Switches can fairly distribute losses using per-VC queueing or FBA
- ❑ UBR: Switches probably will not have separate UBR queues ⇒ Low Fairness unless FBA

## **Bursty UDP Data Traffic**

- □ Metric: Throughput or Efficiency
- Several Client-Server transaction applications use UDP.
- $\Box \text{ Data} \Rightarrow \text{Loss Sensitive} \Rightarrow \text{Retransmission}$
- □ UDP  $\Rightarrow$  No Slow Start  $\Rightarrow$  Losses can continue  $\Rightarrow$  Losses are more expensive than in TCP
- Other conclusions are similar to TCP

#### **Loss-tolerant UDP Traffic**

- Example: Voice over IP
- □ Loss-tolerant generally implies delay sensitive
- $\Box \text{ ATM backbone} \Rightarrow \text{Aggregated flows}$
- ABR: Queues in the router. If hierarchically coded and drop preference indication in packets
   ⇒ Routers can drop the low priority packets
- CBR: Low efficiency due to traffic variability.
   But Routers can drop the low priority packets.
- GFR/VBR/UBR: Packets may enter ATM network and dropped there. CLP bit coded by drop preference.

### **Differentiated Services**

- Details of DS are yet to be finalized.
- Currently 4 queues and 3 drop preferences (July IETF Meeting)
- □ ATM has only two drop preferences: CLP = 0 or 1
- ❑ ABR: Queues in the Router ⇒ Routers can set different thresholds for different drop preferences
- CBR: Queues in the router.
   But not as efficient as ABR for Bursty traffic.
- □ GFR/VBR/UBR: Queues in side the network
   ⇒ Can't handle more than 2 drop preferences

## **Differentiated Services - Priorities**

- □ Four Queues: With Priority and weights
- $\Box Weights \Rightarrow Guaranteed bandwidth$
- ABR/CBR: All queues in the routers
   ⇒ Edge routers can keep multiple priority queues feeding to a single ABR VC
- □ GFR/VBR/UBR: No queues in the routers
   ⇒ Can't enforce priorities in the router
- □ GFR: Higher MCR  $\neq$  Higher Priority
  - $\neq$  Higher share of extra bandwidth
- □ VBR: Higher SCR/PCR  $\neq$  Higher Priority The Ohio State University



- □ ABR: Key Distinction is feedback
   ⇒ Network is congestion free and maximally utilized
- ABR gives more control to edge-routers.
   Routers have more control over drop policies
- □ Other services depend more upon ATM switches ⇒ Fairness difficult to achieve if one VC contains multiple TCP flows The Ohio State University TCP flows Raj Jain

# **Summary (Cont)**

- With ABR it is possible to make use of added buffering in the routers
- □ For Bursty Data: ABR > GFR > VBR > UBR > CBR
- Because of implementation complexity GFR may dominate in the short term
- With ABR, it is possible to implement multiple hierarchical levels of coding
   ⇒ Possible to allow multiple drop preferences
- All other classes can't handle more than two levels of drop preferences ⇒ ABR may rebound if multiple drop preferences in Differentiated Services

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## **Summary (Cont)**

Large careers need ABR to keep queues manageable in the network

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